SHA3-based MACs

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SHA-3 2014 Workshop
August 22, 2014

Outline

- Predecessors
 - FIPS 198 (HMAC)
 - SP 800-108 (KDFs from PRFs)
 - Key Pack
- New specifications
 - KMAC
 - XMAC
 - XKDF
 - Block sizes for SHA3-HMAC

FIPS 198

- Specifies HMAC
 HMAC (K, text)= H(K⊕opad || H(K⊕ipad || text))
- Extra complication needed because of the length extension property of SHA1 and SHA2
- Parameterized by a Block Size
- Allows Truncation

SP 800-108

- Specifies KDF Constructions built from a PRF
 - Counter, Feedback, Double-Pipeline
 - PRF can be (untruncated) HMAC or CMAC
 - Ambiguities can occur if the PRF is variable length (e.g. a MAC based on the SHAKEs)
 - E.g. Counter mode could produce outputs related as A|B|C|D, A|C|E|G
 - We will solve this problem by only considering KMAC, but not XMAC, to be a PRF

Key Pack

- Defined by the Keccak team in "Keyak" and "Ketje" CAESAR entries
- Suggested for KMAC
- Definition:

```
Keypack(K, \ell) = enc<sub>8</sub>(\ell/8) | | K | | 10 \ell-len(K)-9
```

- ℓ is the length of the whole key pack in bits.
- ℓ must be a multiple of 8 between len(K) + 9 and 255*8.
- Need ℓ to be defined further for compatibility.

KMAC

- Usable as a MAC or PRF
- Defined in terms of drop ins.
- Definition:

```
MAC(text) = KMAC(K, text) = H(Keypack(K, \ell) | | text)
```

• $\ell = 8*[(len(K)+9)/8]$

XMAC

- Usable as a MAC, but not a PRF
- Defined in terms of XOFs
- Definition:
 MAC(text) = XMAC(K, text, λ) = X(Keypack(K, ℓ) | | text, λ)
- $\ell = 8*[(len(K)+9)/8]$
- Choice of λ is based on FIPS 198 rules for truncation of HMAC (at least 32 bits.)

XKDF

 Since SP 800-108 KDFs can't use XMAC, we define a KDF that can use XMAC.

• Definition:

 $K_0 := XMAC(K_1, Label || 0x00 || Context || [L]_2, L)$

HMAC block sizes

Algorithm	HMAC Block Size (in Bits)
SHA3-224	1152
SHA3-256	1088
SHA3-384	832
SHA3-512	576

Questions?